

Decision Procedures  
 An Algorithmic Point of View  
 Linear Arithmetic

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Part V

Linear Arithmetic

The Omega Test  
 Outline

- 1 Preprocessing
- 2 Three projections for the Omega Test
- 3 The real shadow
- 4 The dark shadow
- 5 The grey shadow

The Omega Test

- Goal: Decide satisfiability of conjunction of linear constraints over integers

$$\sum_{0 \leq i \leq n} a_i x_i \geq 0$$

- Original application: program optimizations done by a compiler
- Extension of *Fourier-Motzkin* variable elimination:
  - Pick one variable and eliminate it
  - Continue until all variables but one are eliminated

Preprocessing (1)

Normalize coefficients: divide by the GCD

$$\begin{array}{l} 8x + 6y \leq 0 \longrightarrow 4x + 3y \leq 0 \\ 4y \geq 1 \longrightarrow y \geq \lceil 1/4 \rceil \\ 3x + 3y = 2 \longrightarrow x + y = 2/3 \longrightarrow \text{UNSAT} \end{array}$$

'Tightening'

Preprocessing (2)

Eliminate equalities

- Let  $x_i$  denote a variable and  $a_i$  its coefficient, for  $i = 1, 2, \dots$
- Use *substitution* if there is a variable with coefficient  $a_i = 1$
- Otherwise, pick variable  $x_k$  from an equality, and make  $a_k$  positive
- Define  $a \widehat{\text{mod}} b := a - b \lfloor a/b + 1/2 \rfloor$
- Let  $m = a_k + 1$
- Note that  $a_k \widehat{\text{mod}} m = -1$

## Preprocessing (2)

### Eliminate equalities

- Create **new variable**  $\sigma$  and add:

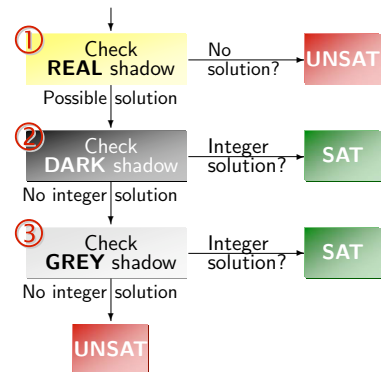
$$m\sigma = \sum_i (a_i \widehat{\text{mod}} m) x_i$$

- Solve for  $x_k$ :

$$x_k = -m\sigma + \sum_{i \neq k} a_k (a_i \widehat{\text{mod}} m) x_i$$

- Q: What is the point of adding a constraint to eliminate one?

## Overview of the Omega Test



## The real shadow

### 1 Check REAL shadow

- Assume we eliminate variable  $z$
- For each pair of upper/lower bound:

$$\begin{array}{ll} \beta \leq bz & az \leq \alpha \\ a\beta \leq abz & abz \leq b\alpha \end{array}$$

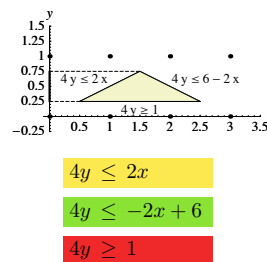
- Constraint for real shadow:

$$a\beta \leq b\alpha$$

- Add this constraint, and call Omega recursively!

## The real shadow: Example I

### 1 Check REAL shadow



Eliminate  $x$ :

$$\begin{array}{ll} 4y \leq 2x & 4y \leq -2x + 6 \\ \hline & 4y \leq 6 - 4y \\ & 8y \leq 6 \end{array}$$

Real Shadow:

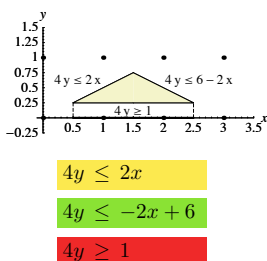
$$\begin{array}{ll} 8y \leq 6 & \Rightarrow y \leq 0.75 \\ 4y \geq 1 & \Rightarrow y \geq 0.25 \end{array}$$

No integer solution

$\Rightarrow$  Original problem has no solution

## The real shadow: Example II

### 1 Check REAL shadow



Let's eliminate  $y$  instead:

$$\begin{array}{ll} 1 \leq 4y & 4y \leq 2x \\ \hline & 1 \leq 2x \\ 1 \leq 4y & 4y \leq -2x + 6 \\ \hline & 1 \leq -2x + 6 \end{array}$$

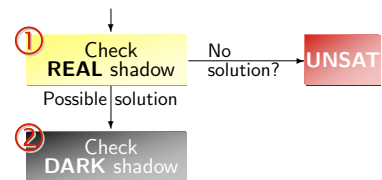
Real Shadow:

$$\begin{array}{ll} 1 \leq 2x & \Rightarrow x \leq 0.5 \\ 1 \leq -2x + 6 & \Rightarrow x \geq 2.5 \end{array}$$

Integer solution!

But original problem has no integer solution!

## From Real to Dark Shadow



- An integer solution for the REAL shadow **does not guarantee** that there is an integer solution for the original problem
- Thus, we check the **DARK shadow** next

## Idea of the dark shadow

### 2 Check DARK shadow

- Idea of the DARK shadow:

$$\begin{array}{l} \beta \leq bz \quad | : b \\ \frac{\beta}{b} \leq z \end{array} \quad \begin{array}{l} az \leq \alpha \quad | : a \\ z \leq \frac{\alpha}{a} \end{array} \quad z \in \mathbb{N}$$

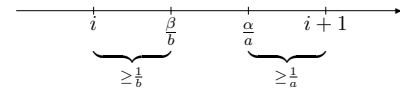
- How to compute the dark shadow?
- Try to *prove* that there is an integer  $z$  between  $\frac{\beta}{b}$  and  $\frac{\alpha}{a}$

## Dark shadow: Proof by contradiction

### 2 Check DARK shadow

Assume there is no integer  $z$  between  $\frac{\beta}{b}$  and  $\frac{\alpha}{a}$ . Then:

Let  $i := \lfloor \frac{\beta}{b} \rfloor \quad i \in \mathbb{Z}$



## Dark shadow: Proof by contradiction

$$\begin{array}{l} \frac{\beta}{b} - i \geq \frac{1}{b} \\ i + 1 - \frac{\alpha}{a} \geq \frac{1}{a} \\ \hline \frac{\beta}{b} + 1 - \frac{\alpha}{a} \geq \frac{1}{b} + \frac{1}{a} \quad | \cdot a \cdot b \\ a\beta + ab - b\alpha \geq a + b \quad | - ab \\ a\beta - b\alpha \geq -ab + a + b \quad | \cdot (-1) \\ \color{red}{b\alpha - a\beta \leq ab - a - b} \end{array}$$

## Dark shadow: Proof by contradiction

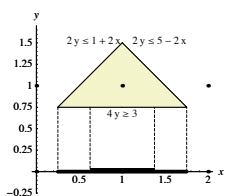
- From previous slide:

$$\begin{array}{l} b\alpha - a\beta \leq ab - a - b \\ \Leftrightarrow \neg(b\alpha - a\beta > ab - a - b) \\ \Leftrightarrow \neg(b\alpha - a\beta \geq ab - a - b + 1) \\ \Leftrightarrow \neg(b\alpha - a\beta \geq (a-1)(b-1)) \end{array} \quad *$$

- Thus, if  $*$  holds, we know that there must be an integer solution.
- If  $a = 1$  or  $b = 1$ , then this is the same as the real shadow. This case is called an **exact projection**.

## Example for the dark shadow

### 2 Check DARK shadow



$$\begin{array}{l} 2y \leq 2x + 1 \\ 2y \leq -2x + 5 \\ 4y \geq 3 \end{array}$$

Eliminate  $y$  with the dark shadow:

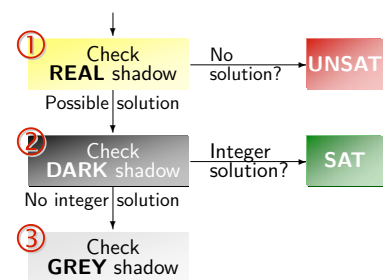
$$\begin{array}{l} 2y \leq 2x + 1 \quad 4y \geq 3 \\ 4(2x + 1) - 2 \cdot 3 \geq (2-1)(4-1) \\ 4y \geq 3 \quad 2y \leq -2x + 5 \\ 4(-2x + 5) - 2 \cdot 3 \geq (2-1)(4-1) \end{array}$$

Dark Shadow:

$$\begin{array}{l} x \geq 5/8 \\ x \leq 11/8 \end{array}$$

$\Rightarrow$  Integer solution!

## From Dark to Grey Shadow



- No integer solution in the DARK shadow **does not guarantee** that there is no integer solution for the original problem
- Thus, we check the **GREY shadow** next

## The grey shadow

### ③ Check GREY shadow

#### Idea of the Grey shadow

If the real shadow  $R$  has integer solutions, but the dark shadow  $D$  does not, search  $R \setminus D$ .

$$\begin{aligned} \text{In } R: & \quad b\alpha \geq abz \geq a\beta \\ \text{Not in } D: & \quad ab - a - b \geq b\alpha - a\beta \\ & \Leftrightarrow ab - a - b + a\beta \geq b\alpha \\ & \Rightarrow ab - a - b + a\beta \geq abz \geq a\beta \quad | : a \\ & \quad (ab - a - b)/a + \beta \geq bz \geq \beta \end{aligned}$$

## The grey shadow

### ③ Check GREY shadow

- Try all values of  $z$  such that

$$(ab - a - b)/a + \beta \geq bz \geq \beta$$

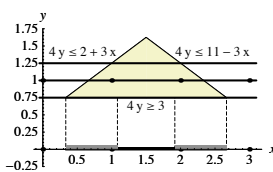
- Optimization: find the largest coefficient  $a$  in any upper bound and try the following for each lower bound  $bz \geq \beta$ :

$$bz = \beta + i \quad \text{for } (ab - a - b)/a \geq i \geq 0$$

- As before, combine this with the original problem, and solve recursively.

## Example of the grey shadow

### ③ Check GREY shadow



$$\begin{aligned} 4y &\leq 3x + 2 \\ 4y &\leq -3x + 11 \\ 4y &\geq 3 \end{aligned}$$

- Eliminate  $y$ :  
 $a = 4, b = 4, \beta = 3$
  - New constraint:  
 $4y = 3 + i$  for  
 $2 \geq i \geq 0$ :  
 $4y = 3$   
 $4y = 4$   
 $4y = 5$
- $\Rightarrow$  Integer solution with  $4y = 4$